

# Lecture 5: Quantum Channel Capacity

## Some Quantum Channels

Q-channel capacity:

$$Q(Q) := \lim_{\varepsilon \rightarrow 0} \limsup_{n \rightarrow \infty} \left\{ \frac{k}{n} : \exists K, C, D \cdot F_p(D Q^{\otimes n} C) > 1 - \varepsilon \right\}$$

Variations:

- quantum channel capacity with two-way classical communication  $Q_2(Q)$ 
  - classical communication does not allow to transmit quantum states

- classical capacity of a quantum channel

$$\mathcal{K}(Q) := \lim_{\varepsilon \rightarrow 0} \limsup_{n \rightarrow \infty} \left\{ \frac{k}{n} : \exists k, C, D \forall |l_k\rangle \in \{|0\rangle, |1\rangle\}^{\otimes k} \quad F_p(|l_k\rangle, D Q^{\otimes n} C) > 1 - \varepsilon \right\}$$

$$Q(Q) \leq \mathcal{K}(Q)$$

- entanglement assisted capacities:  
sender and receiver share some maximally entangled states

examples:

- teleportation.

send a quantum state using only classical information and shared entanglement

$$\text{resources: } 1 \text{ ebit} + 2 \text{ bit} \stackrel{?}{=} 1 \text{ qbit}$$

$\uparrow$                      $\uparrow$                      $\uparrow$   
 1 EPR-pair      2 classical bits      1 qbit

- dense coding:

- send one half of an EPR pair to Bob
- Alice performs one out four operations on her qbit

$$U = O_x^{b_0} O_z^{b_1}, b_0, b_1 \in \{0, 1\}$$

- Alice sends her qbit to Bob
- Bob performs a measurement in the Bell basis  $\Rightarrow$  four different outcomes

$$1 \text{ ebit} + 1 \text{ qbit} \stackrel{?}{=} 2 \text{ bit}$$

# Some particular quantum channels

## 5.1 Depolarizing channel

The quantum state  $\rho$  is transmitted faithfully with probability  $1-p$  or replaced by a random state with probability  $p$ .

$$\rho \mapsto D_p^{\text{pol}}(\rho) = (1-p)\cdot \rho + p \cdot \frac{\mathbb{I}}{\text{tr}(\mathbb{I})}$$

Warning: The completely mixed state  $\frac{\mathbb{I}}{\text{tr}(\mathbb{I})}$  does have a non-zero overlap with any state.

Error operators for  $\mathcal{H} = \mathbb{C}^2$ :

$$D_p^{\text{pol}}(\rho) = (1 - \frac{3}{4}p) \cdot \rho + \frac{p}{4} \left( G_x \rho G_x + G_y \rho G_y + G_z \rho G_z \right)$$

$\Rightarrow$  error operators

$$\left\{ \sqrt{1-\frac{3}{4}p}, \frac{G_x \sqrt{p}}{2}, \frac{G_y \sqrt{p}}{2}, \frac{G_z \sqrt{p}}{2} \right\}$$

## 5.2 Generalized Qudit-Depolarization Channel

(4)

error operators  $\{\mathbb{I}, \sigma_x, \sigma_y, \sigma_z\}$  with different weights, i.e.

$$\mathcal{D}_{(p_x, p_y, p_z)}(\rho) = (1 - p_x - p_y - p_z) \cdot \rho + p_x \sigma_x \rho \sigma_x + p_y \cdot \sigma_y \rho \sigma_y + p_z \sigma_z \rho \sigma_z$$

## 5.3 Dephasing Channel

Scenario: Assume that the channel measures the state with respect to an OVB  $\mathcal{B} = \{|b_i\rangle : i \in I\}$

$\Rightarrow$  The pure input state  $|y\rangle = \sum_{i \in I} \alpha_i |b_i\rangle$  is transformed into the mixture

$$\rho = \sum_{i \in I} |\alpha_i|^2 |b_i\rangle \langle b_i|$$

channel:

$$\rho \rightarrow \mathcal{D}_{p, \mathcal{B}}^{\text{ph}}(\rho) := (1-p) \cdot \rho + p \sum_{i \in I} |b_i\rangle \langle b_i| \rho |b_i\rangle \langle b_i|$$

note:  $|0\rangle \langle 0| = \frac{1}{2} (\mathbb{I} + \sigma_z)$ ,  $|1\rangle \langle 1| = \frac{1}{2} (\mathbb{I} - \sigma_z)$

## 5.4 Results on the channel capacity

Even for the depolarizing which has a lot of symmetry and only one parameter, the channel capacity is not known (in general).

Results (from Bennett, DiVincenzo, Shor, Smolin etc) in 1997/98

1. For  $p < 0,25408$  all of  $Q$ ,  $Q_2$ , and  $\mathcal{K}$  are positive.
2.  $1/3 < p < 2/3$ , the quantum capacity  $Q = 0$ , but both  $Q_2$  and  $\mathcal{K}$  are positive.
3.  $2/3 \leq p < 1$ , both  $Q$  and  $Q_2$  are zero, but  $\mathcal{K}$  is positive.
4. For  $p = 1$ ,  $Q = Q_2 = \mathcal{K} = 0$ .

# The classical capacity of the depolarizing channel

(Ch. King, IEEE Trans. on Information Theory, 2003)

Scenario: Use quantum states  $\rho_i$  to send classical information over the quantum channel

$$D_p^{\text{out}}(\rho) = (1-p)\rho + p \cdot \frac{I}{\text{tr}(\rho)}$$

Input alphabet  $X$  with some distribution  
 $\Rightarrow$  ensemble  $\Sigma$  of quantum states

Output alphabet  $Y$  corresponds to some measurements  $\mathcal{M}$  (or POVM) on the output.

Single use capacity:

$$\mathcal{K}(Q)^{\text{single use}} = \sup_{\Sigma, \mathcal{M}} I(X; Y)$$

Capacity:

$$\mathcal{K}(Q) = \limsup_{n \rightarrow \infty} \frac{1}{n} \sup_{\Sigma, \mathcal{M} \in X^{\otimes n}} I(X; Y)$$

for some channels:

$$\frac{1}{2} \sup_{\Sigma, \mathcal{M} \in \mathcal{X}^{\otimes 2}} I(X^{(2)}; Y^{(2)}) > \sup_{\Sigma, \mathcal{M} \in \mathcal{X}} I(X^u; Y^u)$$

Super additivity

(7)

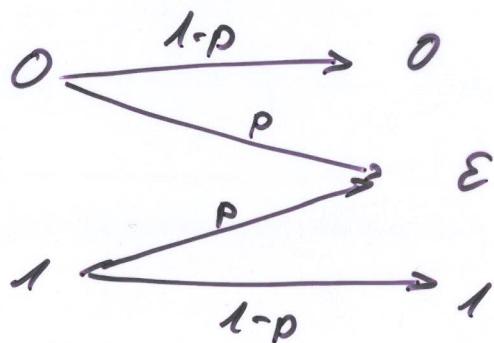
but: For the depolarizing channel, the "classical capacity" is additive, i.e. we don't need to take the limit  $n \rightarrow \infty$ .

Moreover, the capacity is achieved using some orthogonal pure states as inputs and measure in the same basis.

additivity: for  $n > 1$ , use ~~the~~ product states and individual (product) measurements

## 5.5 Quantum erasure channel

recall: BEC is given by



Input:  $\dim \mathcal{H}_m = d$

Output:  $\dim \mathcal{H}_{\text{out}} = d+1$  using an additional orthogonal state  $(\varepsilon)$

$$QEC_p(g) = (1-p) \cdot g + p \cdot |\varepsilon\rangle\langle\varepsilon|$$

## capacities of the QEC:

$$Q(QEC_p) = \max \{0, 1 - 2p\}$$

$$\mathcal{R}(QEC_p) = Q_2(QEC_p) = 1 - p$$

Proof: see Bennett, DiVincenzo & Smolin 97

### conditions:

(a) (b) phase erasure channel:

input: one qubit

output: two qubits, w/o one qubit to indicate the error event

$$s \mapsto (1-p)s \otimes |0\rangle\langle 0|$$

$$+ p\left(\frac{1}{2}s + \frac{1}{2}\sigma_z s \sigma_z\right) \otimes |1\rangle\langle 1|$$

combination of "full erasure" and "phase erasure"

$$s \mapsto (1-\delta-\varepsilon)s \otimes |0\rangle\langle 0|$$

$$+ \varepsilon |1\rangle\langle 1| \otimes |0\rangle\langle 0|$$

$$+ \delta\left(\frac{1}{2}s + \frac{1}{2}\sigma_z s \sigma_z\right) \otimes |1\rangle\langle 1|$$

capacities:

$$Q = \max \{0, 1 - \delta - 2\varepsilon\}$$

$$Q_2 = 1 - \delta - \varepsilon$$

$$K = 1 - \varepsilon$$

### Other types of channels

- channels arising from open quantum, including non-unitary dynamics e.g. by monitoring a system for quantum jumps
- collective effects, e.g. apply the same (local) transformation to all subsystems

$$S \mapsto (1 - p_x - p_y - p_z)S + \sum_{i=x,y,z} p_i \sigma_x^{\otimes n} S \sigma_i^{\otimes n}$$

interchanging particles with some probability (SWAP)

$$\text{SWAP} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} = \frac{1}{2} (\mathbb{I} \otimes \mathbb{I} + \sigma_x \otimes \sigma_x + \sigma_y \otimes \sigma_y + \sigma_z \otimes \sigma_z)$$