

Lecture 18. (Still) Encoding Circuits

18.1 Stabilizer States / Graph States

stabilizer S with $|S| = 2^n$

\Rightarrow unique +1-eigenstate, code $C = [I_n, O, dI]$

stabilizer matrix $S = [X \mid Z]$ with $X, Z \in \mathbb{F}_2^{n \times n}$
 $\text{rk}(S) = n$

normal form for computing an encoding circuit: local operations

\rightarrow matrix X in triangular form $H, P = \begin{pmatrix} 1 & 0 \\ 0 & I \end{pmatrix}$

here even $X = I$ by row operations

commutation of the stabilizer $X \cdot Z^t - Z \cdot X^t = 0$

$\xrightarrow{X=I} Z^t = Z$ (symmetric matrix)

local operations: diagonal of Z is zero

$\Rightarrow Z$ defines an ^{undirected} graph on n vertices

Example: binary code $C' = [8, 4, 4] = C^\perp$

$$G = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix} \quad G \cdot G^t = 0$$

CSS construction with $C_1 = C_2 = C = C^\perp$

one state: $\frac{1}{4} \sum_{c \in C} |c\rangle \quad (*)$

$$\text{Stabilizer } S_0 = \left[\begin{array}{c|cc} G & 0 \\ \hline 0 & G \end{array} \right] \in \mathbb{F}_2^{8 \times 16}$$

$$S_0 = \left[\begin{array}{cccc|cccc|cccc|cccc} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \end{array} \right] \xrightarrow{\text{row operations}} \left[\begin{array}{cccc|cccc|cccc|cccc} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 \end{array} \right]$$

every column
(pair) contains Hadamard on the last
at least one of positions
non-zero entry
 \rightarrow local diff $\rightarrow G_X$

CSS-code \Rightarrow graph is bi-partite

$$S_1 = \left[\begin{array}{cccc|cccc|cccc|cccc} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 \end{array} \right]$$

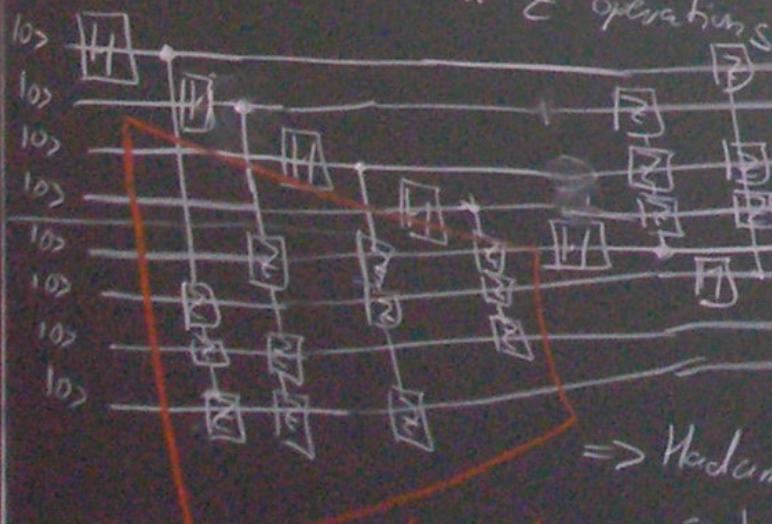
encoding for S_1 ("graph state")

- every 1 on the diagonal of X

- is replaced by II (acting $|0\rangle$)

- followed by controlled operations on the other qubits

here: only controlled-Z operations



encoding for CSS-code

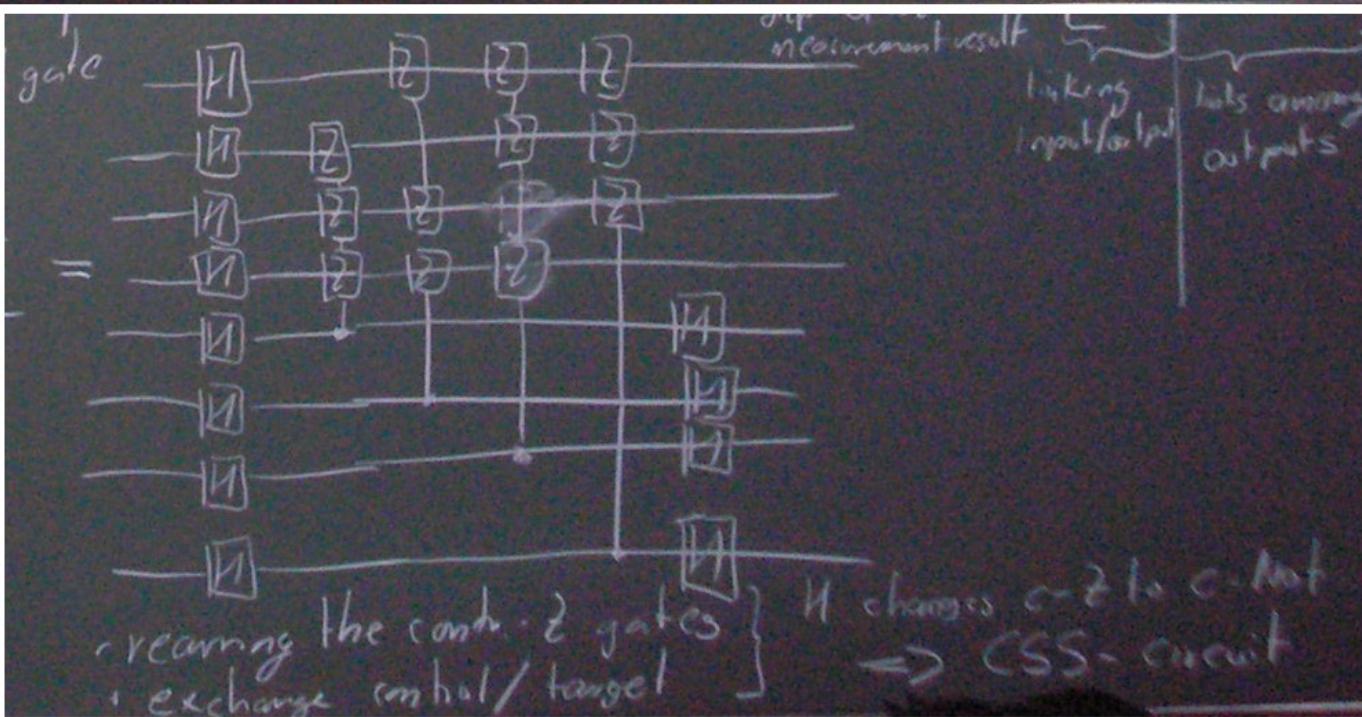
$$|\phi\rangle = \frac{1}{\sqrt{q}} \sum_{c.c. \text{ from } G} |c\rangle$$



circuit for S_1 for S_0

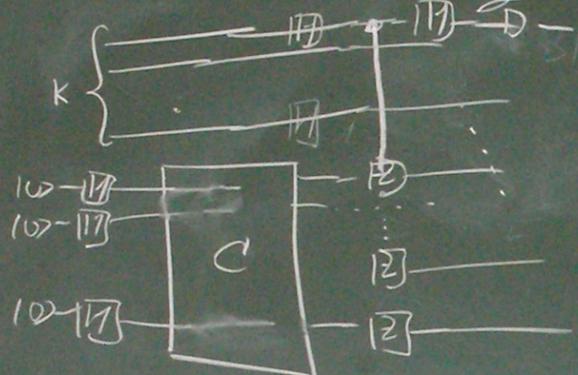
local clifford gate

\Rightarrow He demand on every input
controlled-Z for every edge



outlook graph codes. Graph with K input nodes A
n output nodes B

graph codes. encoding circuit on
K inputs & n outputs
Measuring K qubits



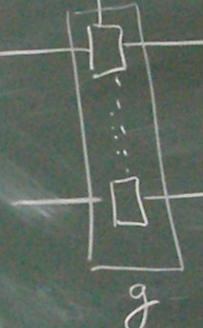
encoding a graph state

$$|\phi_0\rangle = |0 \dots 0\rangle$$

canonical basis

$$|i_1, \dots, i_K\rangle = \overline{X_1^{i_1} \cdot X_K^{i_K}} |0 \dots 0\rangle$$

subcircuit of the form



\Rightarrow projects onto the
+1/-1 eigenspace of g
(phase estimation
algorithm)

References

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